# ON GENERATING SEMIGROUPS FROM A FUNCTION 

LEJO J. MANAVALAN*, P.G. ROMEO<br>Department of Mathematics, Cochin University of Science and Technology, Kalamassery 682022, India

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#### Abstract

A semigroup is a well studied structure so is a function. This paper describes a way in generating a semigroup from a function on a finite set, using its directed graph. Further this paper also studies the properties of the semigroup generated.


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## 1. Introduction

Let $\alpha$ be a function on a set $X$. Recent studies have shown when a function forms an inner translation (also known called Cayley functions) of a semigroup. In 1972 Zupnik characterised all Cayley functions algebraically(in powers of $\beta$ ) [7]. In 2016, Araoujo et all characterised the Cayley functions using functional digraphs[1]. He also described a process(a problem half solved) on generating a semigroup from Cayley functions by constructing a graph with Cayley functions as the vertices and identifying the maximal cliques of the common semigroup graph $G_{S}$ ( two Cayley functions are edge connected if appear on the same semigroup). There are many ways to construct a semigroup from graphs, for example Endomorphism semigroup,

[^0]graph semigroup, commutative graph semigroup, inverse graph semigroup, path semigroup, etc. There is no instance of a semigroup from a function. This paper describes a way in generating a semigroup from a single function on a finite set, using directed graphs, studies the properties of the semigroup generated.

In the sequel $\alpha$ will denote a function mapping a non-empty finite set $X$ onto itself. For any positive integer $n, \alpha^{n}$ denotes the $n^{\text {th }}$ iterate of $\alpha$. By $\alpha^{0}$ we mean the identity function on $S$, so $\alpha^{0}(x)=x$. Let $X$ be a finite set, then $T(X)$ denotes the set of all transformations (functions) from $X$ to $X$.

## 2. Directed Graphs

A directed graph is an ordered pair $D=(S, \rho)$ where the elements of $S$ are called vertices and $A$ is a set of ordered pairs of vertices (binary relation), called arrows, directed edges.

Any pair $(a, b) \in \rho$ is called an arc of $D$, which we will write as $a \rightarrow b$. A vertex $a$ is called is called an initial vertex in $D$ if there is no $b$ in $\rho$ such that $b \rightarrow a$; it is called a terminal vertex in $D$ if there is no $b \in S$ such that $a \rightarrow b$

Definition: A digraph $D$ is called a functional digraph if there is $\alpha: S \longrightarrow S$ such that for all $x, y \in S, x \rightarrow y$ is an arc in $D$ if and only if $\alpha(x)=y$.

Such a functional digraph is denoted as $D_{\alpha}$ as there is one and only one function that represents a functional di-graph.

Let $D$ be a digraph and if there exists pairwise distinct vertices $\ldots, x_{1}, x_{0}, x_{1}, \ldots$ of $D$ such that $x_{0} \longrightarrow x_{1} \longrightarrow \ldots \ldots \longrightarrow x_{k-1} \longrightarrow x_{0}$ then the graph is said to have a cycle of length $k$ denoted as $\left(x_{0} x_{1} \ldots x_{k-1}\right) . D$ has a chain of length $m$, denoted $\left[x_{0} x_{1} \ldots x_{m}\right]$ if there are pairwise distinct vertices such that $x_{0} \longrightarrow x_{1} \longrightarrow \ldots \ldots . \longrightarrow x_{m}$ in $D$. Similarly $D$ has a right ray [left ray or double ray] denoted $\left[x_{0} x_{1} x_{2} \ldots\right\rangle ;\left[\left\langle\ldots x_{2} x_{1} x_{0}\right],\left\langle\ldots x_{1} x_{0} x_{1} \ldots\right\rangle\right]$ if there exist pairwise distinct vertices such that $x_{0} \rightarrow x_{1} \rightarrow x_{2} \rightarrow \ldots \ldots .\left[\left(\ldots \ldots \rightarrow x_{2} \rightarrow x_{1} \rightarrow x_{0}\right)\right.$ or $\left.\left(\ldots \rightarrow x_{-1} \rightarrow x_{0} \rightarrow x_{1} \rightarrow \ldots ..\right)\right]$

Let $S$ be a non-empty set, then $T(S)$ denoted the set of all functions on $S, \alpha \in T(S)$ and $D_{\alpha}$ the directed graph that represents $\alpha$, then a right ray $\left[x_{0} x_{1} x_{2} \ldots.\right\rangle$ in $D_{\alpha}$ is called a maximal right ray if $x_{0}$ is an initial vertex of $D_{\alpha}$.

Definition:Let $D_{\alpha}$ be a functional digraph, where $\alpha \in T(S)$

- A left ray $L=\left\langle\ldots x_{2} x_{1} x_{0}\right.$ ] in $D_{\alpha}$ is called an infinite branch of a cycle $C$ [double ray $W$ ] in $D_{\alpha}$ if $x_{0}$ lies on $C[W]$ and $x_{1}$ does not lie on $C[W]$. We will refer to any such $L$ as an infinite branch in $D_{\alpha}$.
- A chain $P=\left[x_{0} x_{1} \ldots x_{m}\right]$ of length $m \geq 1$ in $D_{\alpha}$ is called a finite branch of a cycle $C$ [double ray $W$, maximal right ray $R$, infinite branch $L$ ] in $D_{\alpha}$ if $x_{0}$ is an initial vertex of $D_{\alpha}, x_{m}$ lies on $C[W, R, L]$ and $x_{m} \geq 1$ does not lie on $C[W, R, L]$. If $x_{m}$ lies on an infinite branch $L=\left\langle\ldots y_{2} y_{1} y_{0}\right]$, we also require that $x_{m} \neq y_{0}$.

The components of $D_{\alpha}$ correspond to the connected components of the underlying undirected graph of $D_{\alpha}$ where the component containing $x$ is defined as follows

Definition: Let $\alpha \in T(S), x \in S$. The subgraph of $D_{\alpha}$ induced by the set

$$
\left\{y \in S: \alpha^{k}(y)=\alpha^{m}(x) \text { for some integers } k, m \geq 0\right\}
$$

is called the component of $D_{\alpha}$ containing $x$.
The following two propositions is due to [4] shows the character of functional digraphs. Proposition: Let $D_{\alpha}$ be a functional digraph. Then for every component $A$ of $D_{\alpha}$, exactly one of the following three conditions holds:
(1) A has a unique cycle but not a double ray or right ray;
(2) A has a double ray but not a cycle ; or
(3) A has a maximal right ray but not a cycle or double ray.

Proposition: Let $D_{\alpha}$ be a functional digraph. Then for every component $A$ of $D_{\alpha}$ :
(1) if A has a (unique) cycle C , then A is the join of C and its branches;
(2) if A has a double ray W , then A is the join of W and its branches;
(3) if A has a maximal right ray R but not a double ray, then A is the join of R and its (finite) branches(type rro).

As we are dealing only with finite functions every component will be a join of a cycle $C$ and its branches.

## 3. Constructing a semigroup from a function

Now consider a function $\alpha \in T(X)$, let $A_{1}, \ldots . . A_{t}$ be the connected components of $D_{\alpha}$ now fix a relation $\leq$ on $A_{1}, \ldots . . A_{t}$ such that $\leq$ is a total ordering on $A_{1}, \ldots \ldots . A_{t}$. Without loss of generality, let $A_{1} \leq \ldots . \leq A_{t}$ (otherwise renumbering as required). Now for each $A_{i}$ fix a $a_{i} \in A_{i}$ such that $a_{i}$ is the initial vertex of the longest branch of the cycle (if the cycle has no branches then choose any one vertex in the cycle ), if there two or more branches of the same length choose any one.

For $x, y \in X$ define $N(x, y)=\left\{(r, q): \alpha^{r}(x)=\alpha^{q}(y)\right\}$. If $N(x, y)$ is empty then $\mathrm{x}, \mathrm{y}$ are in different connected components. If $N(x, y)$ is non-empty then $\mathrm{x}, \mathrm{y}$ are in same connected component and in this case we define $\delta_{x y}=\psi_{x y}-\tau_{x y}$ where $\tau_{x y}=\min \{q:(r, q) \in N(x, y)\}$ and $\psi_{x y}=\min \left\{r: \alpha^{r}(x)=\alpha^{\tau_{x y}}(y)\right\}$. If you take two vertices $x, y$ in the directed graph $D_{\alpha}$, then $\tau_{x y}$ is the minimum distance from $y$ to the path travelled by $x$ as $\alpha$ is composed again and again which is zero if y is on the path, that is $\tau_{x y}=0$ if there is an integer $n$ such that $\alpha^{n}(x)=y$, again if $y^{\prime}=\alpha^{\tau_{x} y}(y)$ then $\psi_{x y}$ is the minimum distance from $x$ to $y^{\prime}$.

The following lemma is required to prove the associativity of the semigroup constructed.
Lemma 1: Let $X$ be a finite set and $\alpha \in T(X)$. Let $\delta_{x y}$ be as defined above, then $\delta_{x \alpha^{l}(y)}=\delta_{x y}+l$

Proof. The proof is by induction, when $l=1$ we have to prove that

$$
\delta_{x \alpha(y)}=\delta_{x y}+1
$$

but $\delta_{x y}=\psi_{x y}-\tau_{x y}$ and $\delta_{x \alpha(y)}=\psi_{x \alpha(y)}-\tau_{x \alpha(y)}$ by definition, also $\tau_{x y}=\min \{q:(r, q) \in N(x, y)\}$ and $\psi_{x y}=\min \left\{r: \alpha^{r}(x)=\alpha^{\tau_{x y}}(y)\right\}$

$$
\begin{array}{lr}
\text { if } \tau_{x y}=0, & \tau_{x \alpha(y)}=0 \text { and } \psi_{x \alpha(y)}=\psi_{x y}+1 \\
\text { if } \tau_{x y} \geq 0, & \tau_{x \alpha(y)}=\tau_{x y}+1 \text { and } \psi_{x \alpha(y)}=\psi_{x y}
\end{array}
$$

so either way

$$
\begin{equation*}
\delta_{x \alpha(y)}=\delta_{x y}+1 \tag{1}
\end{equation*}
$$

Now assume that $\delta_{x \alpha^{l}(y)}=\delta_{x y}+l$ to prove $\delta_{x \alpha^{l+1}(y)}=\delta_{x y}+l+1$,

$$
\begin{aligned}
\delta_{x \alpha^{l+1}(y)} & =\delta_{x \alpha^{l}(y)}+1 \text { from (1) } \\
= & \delta_{x y}+l+1 \text { from assumption }
\end{aligned}
$$

Hence proved

Now for two elements $x, y \in X, \alpha \in T(X)$ define
(2) $\quad x * y= \begin{cases}\alpha_{\mid A_{i}}^{\delta_{a_{i}}+1}(y) & \text { if } x \text { and } y \text { are in the same connected component } A_{i} . \\ x & \text { if } x \in A_{i}, y \in A_{j} \text { and } A_{i} \geq A_{j} \\ y & \text { if } x \in A_{i}, y \in A_{j} \text { and } A_{i} \leq A_{j} .\end{cases}$

To show that $(X, *)$ forms a semigroup we need to show that its an associative binary operation $(x * y) * z=x *(y * z) \forall x, y, z \in X$. For this we have five cases,

Case 1: $x, y, z$ are in 5 different components say $A_{i} A_{j} A_{k}$ respectively, then

$$
\begin{aligned}
& x *(y * z)=(x * y) * z=x \text { if } A_{i} \geq A_{j}, A_{i} \geq A_{k} \\
& x *(y * z)=(x * y) * z=y \text { if } A_{j} \geq A_{i}, A_{j} \geq A_{k} \\
& x *(y * z)=(x * y) * z=z \text { if } A_{k} \geq A_{i}, A_{k} \geq A_{j}
\end{aligned}
$$

Case 2: $x, y$ are in one component say $A_{i}$ and $z$ in $A_{j}$ then

$$
\begin{array}{ll}
x *(y * z)=(x * y) * z=x * y & \text { if } A_{i} \geq A_{j} \\
x *(y * z)=(x * y) * z=z & \text { if } A_{i} \leq A_{j}
\end{array}
$$

Case 3: $x, z$ are in one component say $A_{i}$ and $y$ in $A_{j}$

$$
\begin{array}{ll}
x *(y * z)=(x * y) * z=x * z & \text { if } A_{i} \geq A_{j} \\
x *(y * z)=(x * y) * z=y & \text { if } A_{i} \leq A_{j}
\end{array}
$$

Case 4: $y, z$ are in one component say $A_{i}$ and $x$ in $A_{j}$

$$
\begin{array}{ll}
x *(y * z)=(x * y) * z=x & \text { if } A_{i} \geq A_{j} \\
x *(y * z)=(x * y) * z=y * z & \text { if } A_{i} \leq A_{j}
\end{array}
$$

Case 5: $x, y, z$ are in one component say $A_{i}$,
Let $k=\delta_{a_{i} y}+1$ and $l=\delta_{a_{i} x}+1$ then

$$
\begin{aligned}
x *(y * z) & =x * \alpha^{\delta_{a_{i}}+1}(z)=\alpha^{\delta_{a_{i} x}+1}\left(\alpha^{\delta_{a_{i} y}+1}(z)\right)=\alpha^{k+l}(z) \\
(x * y) * z & =\alpha^{\delta_{a_{i}}+1}(y) * z=\alpha^{\delta_{a_{i}\left(\alpha^{\delta_{i} x+1}(y)\right)}}(z)=\alpha^{\delta_{a_{i}\left(\alpha^{l}(y)\right)}+1}(z)=\alpha^{\delta_{a_{i} y}+l+1}(z) \text { by lemma } 1 \\
& =\alpha^{k+l}(z)
\end{aligned}
$$

Hence $*$ is a associative operation on $X$, and so $(X, *)$ is a semigroup, and this summaries to the following theorem,

Theorem 1: Let $X$ be a finite set and $\alpha \in T(X)$ then there exist a semigroup operation on $X$ derived from the function $\alpha$.

So now one can construct a semigroup from any function $\alpha$ on a finite set which I would rather call a functional semigroup and denote by $S_{\alpha}$ the semigroup constructed from a function in the above manner. $S_{\alpha}$ changes with a change in the total order though there is a similarity in them begin that the principal sub-semigroups are isomorphic.
Theorem 2: Let $X$ be a finite set and $\alpha \in T(X), S_{\alpha}^{1} S_{\alpha}^{2}$ be two semigroups constructed from $\alpha$ as described in (2), with two different total ordering but same fixed points, then a principal sub-semigroup of $S_{\alpha}^{1}$ is isomorphic to a principal sub-semigroup of $S_{\alpha}^{2}$.

Proof. Let $x \in X$, be a vertex of $A_{i}$ a connected component of $\alpha$ then $x * x=\alpha^{\delta_{a_{i}}+1}(x)$ in both $S_{\alpha}^{1}$ and $S_{\alpha}^{2}$, invariant of change in the total order, hence the theorem.
$S_{\alpha}$ changes with a change in the choice of the fixed points $a_{i}$ as well, though there is a similarity in them being that the sub-semigroups generated by $a_{i}$, and $b_{i}$, are isomorphic.

Theorem 3: Let $X$ be a finite set and $\alpha \in T(X)$, Suppose that $\alpha$ has t connected components and a connected component (say $A_{k}$ ) such that the length of the longest branch is $s$ and there is at-least two branches of length $s$, let the initial vertex say of these branches be $a_{k}, b_{k}, \ldots$ Let $S_{\alpha}^{1} S_{\alpha}^{2}$ be two semigroups constructed from $\alpha$ one with fixed points $a_{1}, \ldots a_{K}, \ldots, a_{t}$ and the other with fixed points $a_{1}, . . b_{k}, \ldots a_{t}$ with two different choice of fixed points then a principal sub-semigroup generated by $a_{k} S_{\alpha}^{1}$ is isomorphic to a principal sub-semigroup generated by $b_{k}$ of $S_{\alpha}^{2}$. Also the principal sub-semigroup $a_{i}$ is a commutative sub-semigroup.

Proof. Let $a_{i}$ be a fixed point as described, then
$a_{i} * a_{i}=\alpha^{\delta_{a_{i} a_{i}}+1}\left(a_{i}\right)=\alpha\left(a_{i}\right)$,
$a_{i} * \alpha\left(a_{i}\right)=\alpha\left(\alpha\left(a_{i}\right)\right)=\alpha^{2}\left(a_{i}\right)$
$\alpha\left(a_{i}\right) * a_{i}=\alpha^{\delta_{a_{i} \alpha\left(a_{i}\right)}+1}\left(a_{i}\right)=\alpha^{2}\left(a_{i}\right)$ by lemma 1
and hence we get the sub-semigroup generated by $a_{i}=\left\{\alpha^{n}\left(a_{i}\right) ; n \in \mathrm{~N}\right\}$ Also, $\alpha^{n}\left(a_{i}\right) * \alpha^{m}\left(a_{i}\right)=$ $\alpha^{\delta_{a_{i} \alpha^{n}\left(a_{i}\right)}+1}\left(\alpha^{m}\left(a_{i}\right)\right)=\alpha^{\delta_{a_{i} a_{i}}+n+1}\left(\alpha^{m}\left(a_{i}\right)\right)$ by lemma $1=\alpha^{m+n+1}\left(a_{i}\right)=\alpha^{m}\left(a_{i}\right) * \alpha^{n}\left(a_{i}\right)$, hence the commutativity is proved.

To prove that $S_{a_{k}}=\left\{\alpha^{n}\left(a_{k}\right) ; n \in \mathrm{~N}\right\}$ is isomorphic to $S_{b_{k}}=\left\{\alpha^{n}\left(b_{k}\right) ; n \in \mathrm{~N}\right\}$, Let $\phi: S_{a_{k}} \longrightarrow$ $S_{b_{k}}$ be such that $\phi\left(\alpha^{n}\left(a_{k}\right)\right)=\alpha^{n}\left(b_{k}\right)$, then for $x=\alpha^{n}\left(a_{k}\right), y=\alpha^{m}\left(a_{k}\right) \in S_{a_{k}}$

$$
\begin{aligned}
\phi(x * y) & = & \phi\left(\alpha^{n}\left(a_{k}\right) * \alpha^{m}\left(a_{k}\right)\right) & = & \phi\left(\alpha^{m+n+1}\left(a_{k}\right)\right) & = \\
\phi(x) * \phi(y) & = & \phi\left(\alpha^{n}\left(a_{k}\right)\right) * \phi\left(\alpha^{m+n+1}\left(a_{k}\right)\right) & = & \alpha^{n}\left(b_{k}\right) * \alpha^{m}\left(b_{k}\right) & =
\end{aligned} \alpha^{m+n+1}\left(b_{k}\right)
$$

Hence $\phi$ is an isomorphism

Hence if the directed graph $D_{\alpha}$ has at-most one branch in each connected component then the semigroup generated from $\alpha$ is commutative. Moreover if $\alpha$ is a permutation then the semigroup generated is a union of groups. Further for a function $\alpha$, if $D_{\alpha}$ has $t$ connected
components then $S_{\alpha}$ has $t$ maximal-subgroups. Precisely the vertices of a cycle in $D_{\alpha}$ form a group.

Now comparing the number of semigroup on a finite set and the number of function on a finite set one can see that [3]

| n | functions | semigroups | commutative semigroup |
| :---: | :---: | :---: | :---: |
| 2 | 4 | 4 | 3 |
| 3 | 27 | 18 | 12 |
| 4 | 256 | 126 | 58 |
| 5 | 3,125 | 1,160 | 325 |
| 6 | 46,656 | 15,973 | 2,143 |
| 7 | 823,543 | 836,021 | 1,7291 |
| 8 | $16,777,216$ | $1,843,120,128$ | 221,805 |

So now one concludes that most of the semigroups constructed from a function in a manner described above must be isomorphic.

Theorem 4: Let $X$ be a finite set and $\alpha$ and $\beta$ two functions on $X$, then for any $S_{\alpha}$ there is a ordering on the directed graph $D_{\beta}$ and a choice of vertices $b_{i}$ such that $S_{\alpha}$ and $S_{\beta}$ are isomorphic if and only if the directed graphs $D_{\alpha}$ and $D_{\beta}$ are digraph isomorphic.

Proof. Suppose that $D_{\alpha}$ and $D_{\beta}$ are digraph isomorphic with isomorphism $\phi_{D}$ and $S_{\alpha}$ is constructed from $\alpha$ such that $A_{1} \leq \ldots . \leq A_{t}$ and for each $A_{i}, a_{i} \in A_{i}$ is the chosen fixed vertex, then let $b_{i}=\phi_{D}\left(a_{i}\right)$ be the fixed vertices and $B_{i}$ be the connected component such that $b_{i} \in B_{i}$. Then the semigroup constructed with the order $B_{1} \leq \ldots . \leq B_{t}$ and $b_{i}$ the fixed vertices is isomorphic to $S_{\alpha}$ as when $a, b \in A_{i}, \phi_{D}(a * b)=\phi_{D}\left(\alpha^{\delta_{a_{i}}+1}(b)\right)=\beta^{\delta_{\phi_{D}\left(a_{i}\right) \phi_{D}(a)}+1}\left(\phi_{D}(b)\right)($ by digraph isomorphism $)=$ $\phi_{D}(a) * \phi_{D}(b)$, when $a, b$ belong to different components its trivial .

Conversely suppose that $S_{\alpha}$ and $S_{\beta}$ are isomorphic and that $\phi_{S}$ is the semigroup-isomorphism, then $\phi_{S}$ is also the digraph isomorphism for when $\alpha(a)=b$ then by semigroup isomorphism $\phi_{S}\left(a_{i}\right) * \phi_{S}(a)=\phi_{S}\left(a_{i} * a\right)$

$$
\begin{aligned}
\phi_{S}\left(a_{i} * a\right) & =\phi_{S}(\alpha(a))=\phi_{S}(b) \\
\phi_{S}\left(\left(a_{i}\right)\right) * \phi_{S}(a) & =\beta^{\delta_{b_{i} \phi_{S}\left(a_{i}\right)}+1}\left(\phi_{S}(a)\right)=\quad \beta^{\delta_{b_{i} b_{i}}+1}\left(\phi_{S}(a)\right)=\beta\left(\phi_{s}(a)\right)
\end{aligned}
$$

which shows that if there is an edge from $a$ to $b$ in $D_{\alpha}(\alpha(a)=b)$ then $\beta\left(\phi_{S}(a)\right)=\phi_{S}(b)$ so there is an edge from $\phi_{S}(a)$ to $\phi_{S}(b)$ and vice-versa. Hence $D_{\alpha}$ and $D_{\beta}$ are digraph isomorphic.

Remark: If in equation (2) $\alpha_{\mid A_{i}}^{\delta_{a_{i}}+1}(y)$ is replaced by $\alpha_{\mid A_{i}}^{\delta_{a_{I} x}}(y)$ then the isomorphism in theorem 4 is not obtained.

This construction could be extended to the case when $X$ is countably infinite

## Conflict of Interests

The authors declare that there is no conflict of interests.

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[^0]:    *Corresponding author
    E-mail address: lejojmanavalan171@gmail.com
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