

NUMBER OF SPANNING TREES OF NEW JOIN GRAPHS

S. N. DAOUD^{1,2}

¹Department of Applied Mathematics, Faculty of Applied Science, Taibah University, Al-Madinah, K.S.A.

²Department of Mathematics, Faculty of Science, El-Minufiya University, Shebeen El-Kom, Egypt

Abstract. In mathematics, one always tries to get new structures from given ones. This also applies to the realm of graphs, where one can generate many new graphs from a given set of graphs. In this paper we derive simple formulas of the complexity, number of spanning trees, of some new join graphs, using linear algebra, Chebyshev polynomials and matrix analysis techniques.

Keywords: Number of spanning trees; Chebyshev polynomials; Join graphs.

Mathematics Subject Classification: 05C05, 05C50.

1. Introduction

In this workwe deal with simple and finite undirected graphs G = (V, E), where V is the vertex setand E is the edgeset. For a graph G, a spanning tree in G is a tree which has the same vertex set as G. The number of spanning trees in G, also called, the complexity of the graph, denoted by $\tau(G)$, is a well-studied quantity (for long time). A classical result of Kirchhoff, [19] can be used to determine the number of spanning trees for G = (V, E). Let $V = \{v_1, v_2, ..., v_n\}$, then the Kirchhoff matrix H defined as $n \times n$ characteristic matrix H = D - A, where D is the diagonal matrix of the degrees of G and A is the adjacency matrix of G, $H = [a_{ij}]$ defined as follows: (i) $a_{ij} = -1v_i$ and v_j are adjacent and $i \neq j$, (ii) a_{ij} equals the degree of vertex v_i if i = j, and (iii) $a_{ij} = 0$ otherwise. All of co-factors of H are equal to $\tau(G)$. There are

Received April23, 2013

other methods for calculating $\tau(G)$. Let $\mu_1 \ge \mu_1 \ge \dots \ge \mu_p$ denote the eignvalues of H matrix of a p point graph. Then it is easily shown that $\mu_p = 0$. Furthermore, Kelmans and Chelnokov [18] shown that, $\tau(G) = \frac{1}{p} \prod_{k=1}^{p-1} \mu_k$. The formula for the number of spanning trees in a d-regular graph G can be expressed as $\tau(G) = \frac{1}{p} \prod_{k=1}^{p-1} (d - \mu_k) \quad \text{where} \quad \lambda_0 = d, \lambda_1, \lambda_2, \dots, \lambda_{p-1} \quad \text{are the eigenvalues of the}$ corresponding adjacency matrix of the graph. However, for a few special families of graphs there exists simple formulas that make it much easier to calculate and determine the number of corresponding spanning trees especially when these numbers are very large. One of the first such result is due to Cayley[4] who showed that complete graph on *n* vertices, K_n has n^{n-2} spanning trees that he showed $\tau(K_n) = n^{n-2}, n \ge 2$. Another result, $\tau(K_{p,q}) = p^{q-1}q^{p-1}, p,q \ge 1$, where $K_{p,q}$ is the complete bipartite graph with bipartite sets containing p and q vertices, respectively. It is well known, as in e.g., [5,22]. Another result is due to Sedlacek [23] who derived a formula for the wheel on n+1 vertices, W_{n+1} , he showed that $\tau(W_{n+1}) = (\frac{3+\sqrt{5}}{2})^n + (\frac{3-\sqrt{5}}{2})^n - 2$, for $n \ge 3$. Sedlacek [24] also later derived a formula number of spanning trees in a Mobius ladder, M_n for the $\tau(M_n) = \frac{n}{2}[(2+\sqrt{3})^n + (2-\sqrt{3})^n + 2]$ for $n \ge 2$. Another class of graphs for which an explicit formula has been derived is based on a prism. Boesch, et .al. [2,3]. Daoud, [6-16] later derived formulas for the number of spanning trees for many graphs.

The Lucas numbers are the sequence $\{L_n\}_{n=1}^{\infty}$ defined by the linear recurrence equation $L_n = L_{n-1} + L_{n-2}$ (1) With $L_0 = 2, L_1 = 1$ and $L_2 = 3$. The values of L_n for $n = 1, 2, 3, \dots$ are $1, 3, 4, 7, 11, 18, 29, 47, 76, 123, \dots, [1, 17]$.

The Fibonacci numbers are the sequence $\{F_n\}_{n=1}^{\infty}$ defined by the linear recurrence equation:

$$F_n = F_{n-1} + F_{n-2} \tag{2}$$

With $F_0 = 0$ and $F_1 = F_2 = 1$. The values of F_n for $n = 1, 2, 3, \dots$ are

1, 1, 2, 3, 5, 8, 13, 21,, [21].

Now, we can introduce the following lemma:

Lemma 1.1

 $\tau(G) = \frac{1}{n^2} \det(nI - \overline{D} + \overline{A})$ where \overline{A} , \overline{D} are the adjacency and degree matrices of \overline{G} ,

the complement of G, respectively, and I is the $n \times n$ unit matrix.

The advantageof this formula is to express $\tau(G)$ directly as a determinant rather than in terms of cofactors as in Kirchhoff theorem or eigenvalues as in Kelmans and Chelnokov formula.

2. Chebyshev Polynomial

In this section we introduce some relations concerning Chebyshev polynomials of the first and second kind which we use it in our computations.

We begin from their definitions, Yuanping, et. al. [25].

Let $A_n(x)$ be $n \times n$ matrix such that:

$$A_{n}(x) = \begin{pmatrix} 2x & -1 & 0 & \cdots & 0 \\ -1 & \ddots & \ddots & \ddots & \vdots \\ 0 & \ddots & \ddots & \ddots & 0 \\ \vdots & \ddots & \ddots & \ddots & -1 \\ 0 & \cdots & 0 & -1 & 2x \end{pmatrix},$$
 where all other elements are zeros.

Further we recall that the Chebyshev polynomials of the first kind are defined by:

$$T_n(x) = \cos(n \arccos x) \quad (3)$$

The Chebyshev polynomials of the second kind are defined by

$$U_{n-1}(x) = \frac{1}{n} \frac{d}{dx} T_n(x) = \frac{\sin(n \arccos x)}{\sin(\arccos x)}$$
(4)

It is easily verified that

 $U_{n}(x) - 2xU_{n-1}(x) + U_{n-2}(x) = 0 \quad (5)$

It can then be shown from this recursion that by expanding det $A_n(x)$ one gets

$$U_n(x) = \det(A_n(x)), n \ge 1$$
(6)

Furthermore by using standard methods for solving the recursion (3), one obtains the explicit formula

$$U_{n}(x) = \frac{1}{2\sqrt{x^{2}-1}} [(x + \sqrt{x^{2}-1})^{n+1} - (x - \sqrt{x^{2}-1})^{n+1}], n \ge 1 \quad (7)$$

Where the identity is true for all complex x (except at $x = \pm 1$ where the function can be taken as the limit).

The definition of $U_n(x)$ easily yields its zeros and it can therefore be verified that

$$U_{n-1}(x) = 2^{n-1} \prod_{j=1}^{n-1} (x - \cos \frac{j\pi}{n})$$
(8)

One further notes that

$$U_{n-1}(-x) = (-1)^{n-1}U_{n-1}(x)$$
(9)

These two results yield another formula for $U_n(x)$,

$$U_{n-1}^{2}(x) = 4^{n-1} \prod_{j=1}^{n-1} (x^{2} - \cos^{2} \frac{j\pi}{n})$$
(10)

Finally, simple manipulation of the above formula yields the following, which also will be extremely useful to us latter:

$$U_{n-1}^{2}(\sqrt{\frac{x+2}{4}}) = \prod_{j=1}^{n-1} (x - 2\cos\frac{2j\pi}{n})$$
(11)

Furthermore one can show that

$$U_{n-1}^{2}(x) = \frac{1}{2(1-x^{2})} [1-T_{2n}] = \frac{1}{2(1-x^{2})} [1-T_{n}(2x^{2}-1)], \qquad (12)$$

and

$$T_{n}(x) = \frac{1}{2} \left[(x + \sqrt{x^{2} - 1})^{n} + (x - \sqrt{x^{2} - 1})^{n} \right].$$
(13)

Another interesting fact follows by comparing (7) with the will known closed form formula for the Fibonacci numbers F_n

$$F_n = \frac{1}{\sqrt{5}} \left[\left(\frac{1+\sqrt{5}}{2} \right)^n - \left(\frac{1-\sqrt{5}}{2} \right)^n \right]$$
(14)

namely,

$$U_{n-1}(\frac{3}{2}) = F_{2n} = \frac{1}{\sqrt{5}} \left[\left(\frac{3+\sqrt{5}}{2}\right)^n - \left(\frac{3-\sqrt{5}}{2}\right)^n \right]$$
(15)

Also we can prove that :

$$2T_n\left(\frac{3}{2}\right) = L_{2n} = \left(\frac{3+\sqrt{5}}{2}\right)^n + \left(\frac{3-\sqrt{5}}{2}\right)^n \tag{16}$$

Lemma 2.1 Let $B_n(x)$ be $n \times n$ matrix such that:

$$B_{n}(x) = \begin{pmatrix} x & -1 & 0 & \cdots & \cdots & 0 \\ -1 & 1+x & \ddots & \ddots & \ddots & \vdots \\ 0 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 0 \\ \vdots & \ddots & \ddots & \ddots & 1+x & -1 \\ 0 & \cdots & \cdots & 0 & -1 & x \end{pmatrix}$$

Then

$$\det(B_n(x)) = (x-1)U_{n-1}(\frac{1+x}{2}).$$

Proof:

Straightforward induction using properties of determinants and above mentioned definition.

Lemma 2.2 Let $C_n(x)$ be $n \times n$ matrix, $n \ge 3, x > 2$ such that:

$$C_{n}(x) = \begin{pmatrix} x & 0 & 1 & \cdots & \cdots & 1 \\ 0 & x+1 & \ddots & \ddots & \ddots & \vdots \\ 1 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 \\ \vdots & \ddots & \ddots & \ddots & x+1 & 0 \\ 1 & \cdots & \cdots & 1 & 0 & x \end{pmatrix}$$

Then

$$\det(C_n(x)) = (n+x-2)U_{n-1}(\frac{x}{2})$$

Proof: Straightforward induction using the properties of determinants and lemma 2.1. Lemma 2.3 Let $D_n(x)$ be $n \times n$ matrix, $n \ge 3, x \ge 4$ such that :

$$D_n(x) = \begin{pmatrix} x & 0 & 1 & \cdots & 1 & 0 \\ 0 & \ddots & \ddots & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & 0 \\ 0 & 1 & \cdots & 1 & 0 & x \end{pmatrix}$$

Then: $\det(D_n)$

$$D_n(x)) = \frac{2(x+n-3)}{x-3} [T_n(\frac{x-1}{2}) - 1]$$

Proof: Straightforward induction using properties of determinants and above mentioned definitions of Chebyshev polynomial of the first and second kind.

Lemma 2.4

Let $E_n(x)$ be $n \times n$ matrix, $x \ge 2$ such that:

$$E_{n}(x) = \begin{pmatrix} x & 1 & \cdots & \cdots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 \\ 1 & \cdots & \cdots & \cdots & 1 & x \end{pmatrix}$$

Then

$$\det(E_n) = (x + n - 1)(x - 1)^{n-1}$$

Proof From the definition of the circulant determinants , we have:

$$\det(E_{n}(x)) = \det\begin{pmatrix} x & 1 & \cdots & \cdots & 1 \\ 1 & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & 1 \\ 1 & \cdots & \cdots & 1 & x \end{pmatrix} = \prod_{j=1}^{n} (x + \omega_{j} + \omega_{j}^{2} + \omega_{j}^{3} + \dots + \omega_{j}^{n-1})$$
$$= (x + 1 + 1 + \dots + 1) \times \prod_{j=1,\omega_{j} \neq 1}^{n} (x + \omega_{j} + \omega_{j}^{2} + \omega_{j}^{3} + \dots + \omega_{j}^{n-1})$$
$$= (x + n - 1) \times (x - 1)^{n-1}.$$

Lemma 2.5,[20]

Let
$$A \in F^{n \times n}$$
, $B \in F^{n \times m}$, $C \in F^{m \times n}$ and $D \in F^{m \times m}$ assume that A, D ,

are nonsingular matrices. Then:

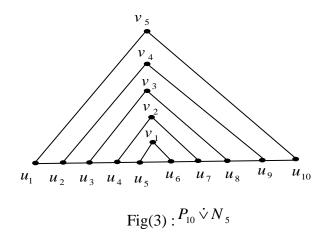
$$\det \begin{pmatrix} A & B \\ C & D \end{pmatrix} = (-1)^{nm} \det(A - BD^{-1}C) \det D = (-1)^{nm} \det A \det(D - CA^{-1}B)$$

Formula in lemma 2.5 give some sort of symmetry in some matrices which facilitates our calculation of determinants.

3. Complexity of some new join Graphs

Definition 3.1

The join graph $P_{2n} \lor N_n$ with 3n vertices, is formed from a simple path P_{2n} with 2n vertices u_1, u_2, \dots, u_{2n} and null graph N_n with n vertices v_1, v_2, \dots, v_n such that $v_n \in N_n$ adjacent with u_1 and u_{2n} in P_{2n} and $v_{n-1} \in N_n$ adjacent with u_2 and u_{2n-1} in P_{2n} and so on. See Fig. (3).



Theorem 3.2 For $n \ge 1$, $\tau(P_{2n} \lor N_n) = 2^{n-1}L_{2n}$.

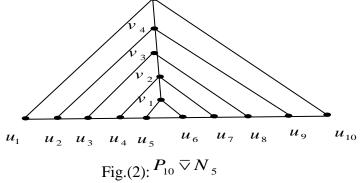
Proof: Applying lemma 1.1, we have:

Straightforward induction using relations between Lucas numbers and properties of Chebyshev polynomials of first kind and second kind, yields:

$$\tau (P_{2n} \circ N_n) = \frac{1}{(3n)^2} [2^{n-1}(n+2)] [\frac{2(3n)^2}{n+2} T_n(\frac{3}{2})] = 2^n T_n(\frac{3}{2})$$
$$= 2^n \times \frac{1}{2} [(\frac{3+\sqrt{5}}{2})^n + (\frac{3-\sqrt{5}}{2})^n] = \frac{1}{2} [(3+\sqrt{5})^n + (3-\sqrt{5})^n] = 2^{n-1} L_{2n}.$$

Definition 3.3

The join graph $P_{2n} \nabla P_n$ with 3n vertices, is formed from a simple path P_{2n} with 2n vertices u_1, u_2, \dots, u_{2n} and path P_n with n vertices v_1, v_2, \dots, v_n such that $v_n \in P_n$ adjacent with u_1 and u_{2n} in P_{2n} and $v_{n-1} \in P_n$ adjacent with u_2 and u_{2n-1} in P_{2n} and so on. See Fig. (2).



Theorem 3.4 Let For $n \ge 2$, $\tau (P_{2n} \bigtriangledown P_n) = 2^{n-1} L_{2n}$.

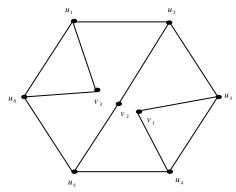
Proof: Applying lemma 1.1, we have:

Straightforward induction using properties of determinants and relations of Chebyshev polynomials of first kind and second kind, yields:

$$\tau (P_{2n} \nabla P_n) = \frac{1}{(3n)^2} [(n+2)U_{n-1}(2)] [\frac{2(3n)^2}{(n+2)} \frac{T_n(\frac{3}{2})U_{n-1}(\frac{5}{2})}{U_{n-1}(2)}] = 2T_n(\frac{3}{2})U_{n-1}(\frac{5}{2})$$
$$= \frac{1}{2^{2n}\sqrt{21}} [(3+\sqrt{5})^n + (3-\sqrt{5})^n] [(5+\sqrt{21})^n - (5-\sqrt{21})^n].$$

Definition 3.5

The join graph $C_{2n} \diamond N_n$ on 3n vertices, is formed from a simple path C_{2n} with 2n vertices u_1, u_2, \dots, u_{2n} and null N_n with n vertices v_1, v_2, \dots, v_n such that $v_n \in N_n$ adjacent with u_1 and u_{2n} in C_{2n} and $v_{n-1} \in N_n$ adjacent with u_2 and u_{2n-1} in C_{2n} and so on. See Fig. (3).



 $\operatorname{Fig}(3): {}^{C_6} \dot{\vee} N_3$

Theorem 3.6: for $n \ge 1$, $\tau(C_{2n} \lor N_n) = 5 \times 2^{n-1} F_{2n}$.

Proof: Applying lemma 1.1, we have:

$$\tau(C_{2n} \lor N_n) = \frac{1}{(3n)^2} \det(3n I - \overline{D} + \overline{A})$$

Straightforward induction using relations between Fibonacci numbers and properties of Chebyshev polynomials of first kind and second kind, yields:

$$\tau(C_{2n} \dot{\vee} N_n) = \frac{\sqrt{5}}{2} [(3 + \sqrt{5})^n - (3 - \sqrt{5})^n] = 5 \times 2^{n-1} F_{2n}$$

Definition 3.7

The join graph $C_{2n} \nabla P_n$ with 3n vertices, is formed from a simple path C_{2n} with 2n vertices u_1, u_2, \dots, u_{2n} and path P_n with n vertices v_1, v_2, \dots, v_n such that $v_n \in P_n$ adjacent with u_1 and u_{2n} in C_{2n} and $v_{n-1} \in P_n$ adjacent with u_2 and u_{2n-1} in C_{2n} and so on. See Fig. (2).

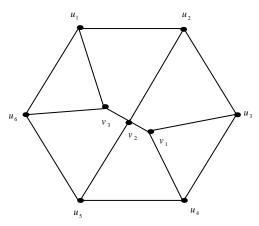


Fig.(4): $C_6 \nabla P_3$

Theorem3.8 For $n \ge 2$,

$$\tau(C_{2n} \nabla P_n) = \frac{5}{2^{2n}\sqrt{105}} [(3+\sqrt{5})^n - (3-\sqrt{5})^n] [(5+\sqrt{21})^n - (5-\sqrt{21})^n]$$

Proof: Applying lemma 3.1, we have:

$$=\frac{1}{(3n)^2}\det C \det(A - BC^{-1}B^T)$$

Straightforward induction using properties of determinants and relations of Chebyshev polynomials of first kind and second kind, yields:

$$\tau(C_{2n} \nabla P_n) = \frac{5}{2^{2n} \sqrt{105}} [(3 + \sqrt{5})^n - (3 - \sqrt{5})^n] [(5 + \sqrt{21})^n - (5 - \sqrt{21})^n]; n \ge 2.$$

4.Conclusion

The number of spanning trees $\tau(G)$ in graphs (networks) is an important invariant. The evaluation of this number is not only interesting from a mathematical (computational) perspective, but also, it is an important measure of reliability of a network and designing electrical circuits. Some computationally hard problems such as the travelling salesman problem can be solved approximately by using spanning trees. Due to the high dependence of the network design and reliability on the graph theory we introduced the above important theorems and lemmas and their proofs.

5. Acknowledgements

The author is deeply indebted and thankful to the deanship of the scientific research for his helpful and for distinct team of employees at Taibah university, Al-Madinah Al-Munawarah, K.S.A. for their continuous helps and the encouragement us to inalize this work. This research work was supported by a grant No. (3021/1434) from the deanship of the scientific research at Taibah university, Al-Madinah Al-Munawwarah, K.S.A.

REFERENCES

- [1] Andrews G., Number theory ,W.B. Saunders Company, London, (1971).
- [2] Boesch, F.T. and Bogdanowicz, Z. R. : The number of spanning trees in a Prism, Inter. J. Comput. Math., 21, 229-243 (1987).
- [3] Boesch, F.T. and Prodinger, H., Spanning tree Formulas and Chebyshev Polynomials, J. of Graphs and Combinatorics 2, 191-200 (1986).
- [4] Cayley, G. A., A Theorm on Trees, Quart. J. Math. 23, 276-378 (1889).
- [5] Clark, L., On the enumeration of multipartite spanning trees of the complete graph, Bull. Of the ICA 38, 50 - 60 (2003).
- [6] Daoud S. N., Some Applications of spanning trees in Complete and Complete bipartite graph : American Journal of Applied Sci. Pub.,9(4) ,584-592, (2012).
- [7] Daoud S. N., Complexity of Some Special named Graphs and Chebyshev polynomials International journal applied Mathematical and Statistics , Ceser Pub. Vol. 32, No. 2,77-84, (2013).

- [8]Daoud S. N. and Elsonbaty A., Complexity of Trapezoidal Graphs with different Triangulations: Journal of combinatorial number theory. Vol.4. no.2, 49-59 (2013).
- [9] Daoud S. N. and Elsonbaty A., Complexity of Some Graphs Generated by Ladder Graph, Journal applied Mathematical and Statistics, Vol. 36, No. 6, 87-94, (2013).
- [10] Daoud S. N., Chebyshev polynomials and spanning tree formulas: International J.Math. Combin, Vol.4, 68-79 (2012).
- [11] Daoud S. N., Number of spanning trees for Splitting of some Graphs, International J. of Math. Sci. & Engg. Appls., Vol. 7, No.II, 169-179, (2013).
- [12]Daoud S. N., Number of spanning trees of Corona of some Special Graphs, International J. of Math. Sci. & Engg. Appls., Vol. 7, No.II, 117-129, (2013).
- [13] Daoud S. N., Number of spanning trees of Join of some Special Graphs, European J. of Scientific Research., Vol. 87, No. 2, 170-181, (2012).
- [14] Daoud S. N., Some Applications of spanning trees of Circulant graphs C_6 and their applications: Journal of Math & Statistics Sci. Pub.,8(1) ,24-31 (2012).
- [15] Daoud S. N., Complexity of Cocktail party and Crown graph : American Journal of Applied Sci. Pub.,9(2) ,202-207 (2012).
- [16] Elsonbaty A. and Daoud S. N.: Number of Spanning Trees of Some Circulant Graphs and Their Asymptotic Behavior: International journal applied Mathematical and Statistics Vol. 30, No. 6, 93-102, (2012).
- [17] Guiduli B., Topics in theory of number, Springer Sciences, Inc, USA (2003).
- [18] Kelmans, A. K. and Chelnokov, V. M.: A certain polynomials of a graph and graphs with an extermal number of trees J. Comb. Theory (B) 16, 197-214 (1974).
- [19] Kirchhoff, G. G., Uber die Auflosung der Gleichungen, auf welche man beiderUntersuchung der LinearenVerteilunggalvanischerStrmegefhrtwird, Ann. Phys. Chem. 72, 497 – 508 (1847).
- [20] Marcus M. A servy of matrix theory and matrix inequalities Unvi.Allyn and Bacon.Inc. Boston (1964).
- [21] Niven I., Introduction to theory of number. John Wiley& Sons, New York (1972).
- [22] Qiao, N. S. and Chen, B. : The number of spanning trees and Chains of graphs, J. Applied Mathematics, no.9, 10-16 (2007).
- [23] Sedlacek, J., Lucas number in graph theory. In: Mathematics (Geometry and Graph theory) (Chech), Univ. Karlova, Prague 111-115 (1970).
- [24] Sedlacek, J., On the Skeleton of a Graph or Digraph. In Combinatorial Structures and their Applications (r. Guy, M. Hanani, N. Saver and J.Schonheim, eds), Gordon and Breach, New York, 387-391(1970).
- [25] Yuanping, Z., Xuerong Y., Mordecai J., Chebyshev polynomials and spanning trees formulas for circulant and related graphs. Discrete Mathematics, 298, 334-364(2005).