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## OPERATORS ON LATTICE-VALUED AUTOMATA

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**Abstract.** The concepts of source core and successor core for residuated lattice valued finite state machine are introduced and their properties are studied. We found the relation between the genetic subset of lattice valued finite state machine and the source core. Using the implication of complete residuated lattice we introduced the fuzzy operators namely fuzzy initial core and fuzzy final core which are generalizations of source core and successor core respectively. Also we have done comparative study of these operators with the source and successor operators.

**Keywords:**  $L$ -valued finite state machine; fuzzy automaton; source, successor and core operators.

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### 1. Introduction

Study of fuzzy automata theory was firstly proposed by Wee in 1967[37] and Wee and Fu in 1969 [38] as a model of learning systems. Since then many authors has been studied fuzzy automaton for characterizations of fuzzy languages [12, 27, 29, 30, 32, 36]. Algebraic study of fuzzy automata theory was started by Malik, Mordeson and Sen [19, 20] and then by many

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others[9, 10, 11, 14, 17, 18, 21]. Topological aspects of fuzzy automata was studied by Das [3], Srivastava and Tiwari [31], Tiwari et.al [35] etc. Many researchers have studied fuzzy automata theory in allied directions to see them we refer to [5, 8, 22, 23].

Generalization of fuzzy automaton as lattice valued automaton was proposed by Qiu [24, 25]. Since then many researchers have studied fuzzy automata with membership values in complete residuated lattice, lattice ordered monoid, orthomodular lattice and other kinds of lattices [2, 4, 13, 15, 16, 26, 33, 34, 39]. Tiwari et.al [33, 34] were studied algebraic and topological characterization of  $L$ -valued finite state machines(called it has also  $L$ -fuzzy automaton) with the underlying structure of lattice ordered monoid.

The paper is arranged in six sections. Preliminary section 2 consist of basic concepts of complete residuated lattice and  $L$ -valued finite state machine. Section 3 contains the concepts of source, successor operators and their fuzzified operators. The operators source core and successor core studied in section 4. Newly introduced operators Fuzzy initial core and fuzzy final core are discussed along with their properties in section 5. Lastly, we conclude the discussion in section 6.

## 2. Preliminaries

In this section, we discuss complete residuated lattice and  $L$ -valued finite state machines. We begin with the definition of residuated lattice.

**Definition 2.1.** [1, 6, 24] *A residuated lattice is a quintuple  $L = (L, \wedge, \vee, *, \rightarrow, 0, 1)$ , where*

- (i)  $(L, \wedge, \vee, 0, 1)$  is a lattice with 0 and 1 as its least and greatest elements respectively;
- (ii)  $(L, *, 1)$  is a commutative monoid with unit 1;
- (iii)  $*$  and  $\rightarrow$  form an adjoint pair, i.e. they satisfy the adjunction property: i.e.  $\alpha * \beta \leq \gamma$  iff  $\alpha \leq \beta \rightarrow \gamma$  for any  $\alpha, \beta, \gamma \in L$ .

In literature the structure of residuated lattice is also called as integral, commutative, residuated,  $l$ -monoid [6]. If in addition,  $(L, \wedge, \vee, 0, 1)$  is a complete lattice, the  $\mathcal{L}$  is called as complete residuated lattice. Throughout this paper,  $L$  stands for complete residuated lattice.

We give an account of some properties of complete residuated lattices as follows: For any

$a, b, x_i \in L$

$$(P_1) (\bigvee_i x_i) * a = \bigvee_i (x_i * a),$$

$$(P_2) a \leq b \Leftrightarrow a \rightarrow b = 1,$$

$$(P_3) a \leq b \Rightarrow c \rightarrow a \leq c \rightarrow b \text{ and } b \rightarrow c \leq a \rightarrow c,$$

$$(P_4) a \rightarrow (\bigwedge_i x_i) = \bigwedge_i (a \rightarrow x_i),$$

$$(P_5) (\bigvee_i x_i) \rightarrow a = \bigwedge_i (a \rightarrow x_i),$$

$$(P_6) 1 \rightarrow a = a,$$

$$(P_7) (a * b) \rightarrow c = a \rightarrow (b \rightarrow c).$$

Additionally, some derived formulae are also necessary for us. We just report them as follows:

i)  $\neg a = a \rightarrow 0, a \in L$

ii) If  $A$  is  $L$ -fuzzy subset of  $Q$  and  $q \in Q$ , then  $Q - A(q) = \neg A(q) = A(q) \rightarrow 0$ .

**Definition 2.2.** A residuated lattice  $\mathcal{L} = (L, \wedge, \vee, *, \rightarrow, 0, 1)$  is said to be a residuated lattice without zero divisors if for all  $a, b \in L$   $a \neq 0, b \neq 0 \Rightarrow a * b \neq 0$ .

**Definition 2.3.** [24] An  $L$ -valued finite state machine (In short,  $L$ -VFSM) or  $L$ -valued semi-automaton is a triple  $\mathcal{M} = (Q, \Sigma, \delta)$ , where  $Q$  and  $\Sigma$  are non empty finite sets of states and input symbols respectively, and  $\delta$  is a  $L$ -fuzzy subset of  $Q \times \Sigma \times Q$  i.e.  $\delta$  is a function from  $Q \times \Sigma \times Q$  into  $L$ , called transition function.

Intuitively,  $\delta(p, \sigma, q)$  stands for the membership degree of the transition from state  $p \in Q$  to the state  $q \in Q$ , when the input is  $\sigma \in \Sigma$ .

**Note 2.1.** The collection of all finite sequence of elements of  $\Sigma$  (called, tapes or string) denoted by  $\Sigma^*$ , and  $\varepsilon$  denotes the empty string in  $\Sigma^*$ . For any  $x \in \Sigma^*$ ,  $|x|$  denotes the length of the string  $x$ . Clearly,  $|\varepsilon| = 0$ .

**Definition 2.4.** [24] Let  $\mathcal{M} = (Q, \Sigma, \delta)$  be  $L$ -VFSM. Extend  $\delta : Q \times \Sigma \times Q \rightarrow L$  to  $\delta^* : Q \times \Sigma^* \times Q \rightarrow L$  inductively as follows:

for any  $p, q \in Q, \sigma \in \Sigma$  and  $x \in \Sigma^*$ ,

$$(1) \delta^*(p, \varepsilon, q) = \begin{cases} 1 & \text{if } p = q \\ 0 & \text{otherwise} \end{cases}$$

$$(2) \delta^*(p, x\sigma, q) = \bigvee_{r \in Q} \{\delta^*(p, x, r) * \delta(r, \sigma, q)\}.$$

Then  $\delta^*$  is called extended transition function with respect to  $*$ .

If we replace  $*$  by  $\wedge$  in the Definition 2.4, then it reduces to.

**Definition 2.5.** Let  $\mathcal{M} = (Q, \Sigma, \delta)$  be  $L$ -VFSM. We extend  $\delta : Q \times \Sigma \times Q \rightarrow L$  to  $\delta^\wedge : Q \times \Sigma^* \times Q \rightarrow L$  as follows:

for any  $p, q \in Q$ ,  $\sigma \in \Sigma$  and  $x \in \Sigma^*$ ,

$$(1) \delta^\wedge(p, \varepsilon, q) = \begin{cases} 1 & \text{if } p = q \\ 0 & \text{otherwise} \end{cases}$$

$$(2) \delta^\wedge(p, x\sigma, q) = \bigvee_{r \in Q} \{\delta^\wedge(p, x, r) \wedge \delta(r, \sigma, q)\}.$$

Then  $\delta^\wedge$  is called extended transition function with respect to  $\wedge$ .

**Theorem 2.1.** [13] Let  $L$  be a lattice and  $X$  be a finite subset of  $L$ . Then  $\wedge$ -semilattice of  $L$  generated by  $X$ ,  $X_\wedge = \{x_1 \wedge \cdots \wedge x_k \mid k \geq 1, x_1, \dots, x_k \in X\} \cup \{1\}$ , is finite. Similarly, the  $\vee$ -semilattice of  $L$  generated by  $X$ ,  $X_\vee = \{x_1 \vee \cdots \vee x_k \mid k \geq 1, x_1, \dots, x_k \in X\} \cup \{0\}$ , is also finite.

The following theorem shows that the importance of complete lattice.

**Theorem 2.2.** Let  $\mathcal{M} = (Q, \Sigma, \delta)$  be  $L$ -VFSM. Then the image set of the extension transition function  $\delta^\wedge : Q \times \Sigma^* \times Q \rightarrow L$  is finite. i.e.  $Im(\delta^\wedge) = \{\delta^\wedge(p, x, q) \mid p, q \in Q, x \in \Sigma^*\}$  is a finite subset of  $L$ .

**Proof.** Let  $X = Im(\delta)$ . As,  $Q$  and  $\Sigma$  are finite, we have  $X$  is finite. We have to prove that  $Im(\delta^\wedge)$  is finite, for this we first show that,  $Im(\delta^\wedge) \subseteq (X_\wedge)_\vee$ .

We prove this, by induction on length of  $x$  denoted by  $|x|$ . When  $|x| = 0$ , then  $\delta^\wedge = 0$  or  $1 \in (X_\wedge)_\vee$  and if  $|x| = 1$ , then  $\delta^\wedge = \delta \in X \subseteq (X_\wedge)_\vee$ .

Now suppose it hold for  $|x| = k$ . We want to prove it for  $|x| = k + 1$ . Let  $x = \sigma_1 \sigma_2 \dots \sigma_k \sigma_{k+1}$ , then we have for any  $p, q \in Q$ ,  $\delta^\wedge(p, \sigma_1 \sigma_2 \dots \sigma_k \sigma_{k+1}, q) = \bigvee_{r \in Q} \{\delta^\wedge(p, \sigma_1 \sigma_2 \dots \sigma_k, r) \wedge \delta(r, \sigma_{k+1}, q)\}$ .

By induction step, we have for any  $r \in Q$ ,  $\delta^\wedge(p, \sigma_1 \sigma_2 \dots \sigma_k, r) \in (X_\wedge)_\vee$  and thus  $\delta^\wedge(p, \sigma_1 \sigma_2 \dots \sigma_k, r) \wedge$

$\delta(r, \sigma_{k+1}, q) \in ((X_\wedge)_\vee)_\wedge = (X_\wedge)_\vee$  and therefore,  $\delta^\wedge(p, \sigma_1 \sigma_2 \dots \sigma_k \sigma_{k+1}, q) \in ((X_\wedge)_\vee)_\vee = (X_\wedge)_\vee$ . Hence  $Im(\delta^\wedge) \subseteq (X_\wedge)_\vee$ . By Theorem 2.1,  $(X_\wedge)_\vee$  is finite subset of  $L$ . Thus  $Im(\delta^\wedge)$ , as a subset of  $(X_\wedge)_\vee$ , is also finite subset of  $L$ .

**Remark 2.1.** *The above Theorem 2.2 shows that, if we use the definition 2.5, we do not need completeness of the lattice, but we need the definition 2.4 have we use complete residuated lattice throughout this paper.*

### 3. Source, successor operators and their fuzzified operators

The study of source and successors operators can be found in Tiwari et.al.[33] and their fuzzified operators in Tiwari et.al.[34], but their underlying lattice structure was lattice ordered monoid. Here, we use complete residuated lattice as the structure of membership values, since structure of residuated lattice,  $(L, *, 1)$  is a one kind of lattice ordered monoid structure. So we collect some of the definitions and results from [33, 34].

**Definition 3.1.** *Let  $\mathcal{M} = (Q, \Sigma, \delta)$  be an  $L$ -valued finite state machine and  $A \subseteq Q$ . The source and the successors of  $A$  are respectively the sets*

$$\sigma(A) = \{p \in Q \mid \delta^*(p, x, q) > 0, \text{ for some } x \in \Sigma^* \text{ and } q \in A\}$$

$$s(A) = \{q \in Q \mid \delta^*(p, x, q) > 0, \text{ for some } x \in \Sigma^* \text{ and } p \in A\}$$

**Definition 3.2.** *An  $L$ -VFSM  $\mathcal{N} = (R, \Sigma, \lambda)$  is called an lattice valued sub finite state machine (In short,  $L$ -VSFSM) of  $L$ -VFSM  $\mathcal{M} = (L, Q, \Sigma, \delta)$ , if  $R \subseteq Q, s(R) = R$  and  $\delta|_{R \times \Sigma \times R} = \lambda$ . Further, this  $L$ -VSFSM is called separated if  $s(Q - R) = Q - R$ .*

**Definition 3.3.** *An  $L$ -VFSM  $\mathcal{M} = (L, Q, \Sigma, \delta)$  is called connected if  $\mathcal{M}$  has no separated proper  $L$ -VSFSM.*

**Definition 3.4.** *Let  $\mathcal{M} = (Q, \Sigma, \delta)$  be an  $L$ -valued finite state machine and  $A \in \mathcal{F}_L(Q)$ . The fuzzy source and the fuzzy successors of  $A$  are respectively defined as:*

$$so(A)(q) = \bigvee_{\substack{x \in \Sigma^* \\ p \in Q}} \{\delta^*(q, x, p) * A(p)\}$$

$$su(A)(q) = \bigvee_{\substack{x \in \Sigma^* \\ p \in Q}} \{A(p) * \delta^*(p, x, q)\}$$

**Theorem 3.1.** Let  $L$  be a complete residuated lattice without zero divisors and  $A \subseteq Q$ . Then for all  $p \in Q$ ,

i)  $so(1_A)(p) > 0$  iff  $p \in \sigma(A)$

ii)  $su(1_A)(p) > 0$  iff  $p \in s(A)$

**Proof.** Proof is same as in [34].

**Definition 3.5.** Let  $\mathcal{M} = (Q, \Sigma, \delta)$  be an  $L$ -VFSM. Then  $R \in \mathcal{F}_L(Q)$  is called an  $L$ -fuzzy lattice valued sub finite state machine (In short,  $L$ -fuzzy  $L$ -VSFSM) of  $\mathcal{M}$ , if  $su(R)(q) \leq R(q), \forall q \in Q$ . Further, this  $L$ -fuzzy  $L$ -VSFSM is called separated, if  $su(Q - R)(q) = Q - R(q), \forall q \in Q$ .

**Definition 3.6.** Let  $\mathcal{M} = (Q, \Sigma, \delta)$  be an  $L$ -VFSM. Then an  $L$ -VFSM  $\mathcal{M}$  is called  $\ell$ -connected, if  $\mathcal{M}$  has no non-constant separated  $L$ -fuzzy  $L$ -VSFSM.

**Theorem 3.2.** Let  $\mathcal{M} = (Q, \Sigma, \delta)$  be an  $L$ -VFSM. If  $\mathcal{M}$  is  $\ell$ -connected, then  $\mathcal{M}$  is connected.

**Proof.** Let  $\mathcal{M}$  be  $\ell$ -connected and  $\mathcal{N} = (R, \Sigma, \lambda)$  be separated  $L$ -VSFSM of  $\mathcal{M}$ . Then  $s(R) = R$  and  $s(Q - R) = Q - R$ .

Now, if  $p \in R$ , then  $1_R(p) = 1$  and  $su(1_R)(p) = \bigvee_{\substack{x \in \Sigma^* \\ q \in Q}} \{1_R(q) * \delta^*(q, x, p) \geq 1_R(p) * \delta^*(p, \epsilon, p) =$

$$1 * 1 = 1 \implies su(1_R)(p) = 1$$

and if  $p \notin R$ , then  $1_R(p) = 0$  and  $su(1_R)(p) = \bigvee_{\substack{x \in \Sigma^* \\ q \in Q}} \{1_R(q) * \delta^*(q, x, p)$

if  $q \notin R$ , then  $1_R(q) = 0$  and if  $q \in R$  and  $\delta^*(q, x, p) > 0$ , then  $p \in s(R)$ , but  $s(R) = R \implies p \in R$ , which is contradiction and thus  $\delta^*(q, x, p) = 0$ . Therefore, for any  $q \in Q$ , either  $1_R(q) = 0$  or  $\delta^*(q, x, p) = 0$  and thus  $su(1_R)(p) = 0$ .

Therefore, for any  $p \in Q$ ,  $su(1_R)(p) = 1_R(p)$ .

Also, since  $Q - 1_R = 1_{Q-R}$ , similarly for any  $p \in Q$ , we have  $su(Q - 1_R)(p) = su(1_{Q-R})(p) = 1_{Q-R}(p) = Q - 1_R(p)$ .

Thus,  $1_R$  is separated  $L$ -fuzzy  $L$ -VSFSM, but since  $\mathcal{M}$  is  $\ell$ -connected, we have  $1_R$  is constant i.e.  $1_R = 1$  or  $0$  i.e.  $R = Q$  or  $\phi$ .

Hence, there does not exist any separated proper  $L$ -VSFSM and thus  $\mathcal{M}$  is connected.

#### 4. Source core and successor core operators

We are chiefly motivated from the concept of core (we say here, source core) defined in the article [33], which is depends on source operator. In this section, we introduce one more core operator, which will depends on successor operator and so we called it successor core.

**Definition 4.1.** *Let  $\mathcal{M} = (Q, \Sigma, \delta)$  be an L-VFSM and  $A \subseteq Q$ . Then source core and successor core of  $A$  are respectively defined as*

$$\begin{aligned}\mu_{\sigma}(A) &= \{q \in A / \sigma(q) \subseteq A\} \text{ and} \\ \mu_s(A) &= \{q \in A / s(q) \subseteq A\}.\end{aligned}$$

We frequently write here,  $\sigma(q)$  and  $s(q)$  for  $\sigma(\{q\})$  and  $s(\{q\})$  respectively. Similarly, we shall frequently write  $\mu_{\sigma}(\{q\})$  and  $\mu_s(\{q\})$  as just  $\mu_{\sigma}(q)$  and  $\mu_s(q)$ .

**Theorem 4.1.** *Let  $\mathcal{M} = (Q, \Sigma, \delta)$  be an L-VFSM. Then for  $A, B \subseteq Q$*

- i) if  $A \subseteq B$ , then  $\mu_{\sigma}(A) \subseteq \mu_{\sigma}(B)$  and  $\mu_s(A) \subseteq \mu_s(B)$*
- ii)  $\mu_{\sigma}(A) \subseteq A$  and  $\mu_s(A) \subseteq A$*
- iii)  $\mu_{\sigma}(A \cap B) = \mu_{\sigma}(A) \cap \mu_{\sigma}(B)$  and  $\mu_s(A \cap B) = \mu_s(A) \cap \mu_s(B)$ .*

**Proof.** We only prove the properties of  $\mu_{\sigma}$ , because the properties of  $\mu_s$  are similar.

(i) Let  $q \in \mu_{\sigma}(A)$ . Then  $\sigma(q) \subseteq A$ . But  $A \subseteq B \implies \sigma(q) \subseteq B \implies q \in \mu_{\sigma}(B)$ . Thus,  $\mu_{\sigma}(A) \subseteq \mu_{\sigma}(B)$ .

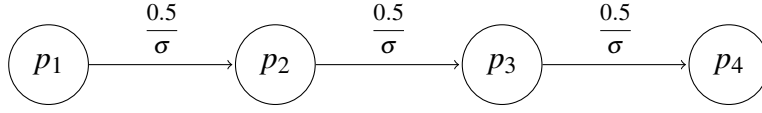
(ii) Let  $q \in \mu_{\sigma}(A)$ . Then  $\sigma(q) \subseteq A$ . But  $q \in \sigma(q)$  always and hence,  $q \in A$ . Thus,  $\mu_{\sigma}(A) \subseteq A$ .

(iii)  $q \in \mu_{\sigma}(A \cap B) \Leftrightarrow \sigma(q) \subseteq A \cap B \Leftrightarrow \sigma(q) \subseteq A$  and  $\sigma(q) \subseteq B \Leftrightarrow q \in \mu_{\sigma}(A)$  and  $q \in \mu_{\sigma}(B) \Leftrightarrow q \in \mu_{\sigma}(A) \cap \mu_{\sigma}(B)$ . Thus,  $\mu_{\sigma}(A \cap B) = \mu_{\sigma}(A) \cap \mu_{\sigma}(B)$ .

**Remark 4.1.** *From, the above theorem, it is obvious that  $\mu_{\sigma}(\mu_{\sigma}(A)) \subseteq \mu_{\sigma}(A)$  and  $\mu_s(\mu_s(A)) \subseteq \mu_s(A)$ . In general  $\mu_{\sigma}(\mu_{\sigma}(A)) \neq \mu_{\sigma}(A)$  and  $\mu_s(\mu_s(A)) \neq \mu_s(A)$ . For this, consider.*

**Example 4.1.** *Let  $L = [0, 1]$  and  $(*, \rightarrow)$  be the Lukasiewicz pair. i.e.  $a * b = \max(0, a + b - 1)$  and  $a \rightarrow b = \min(1, 1 - a + b)$ ,  $a, b \in L$ . Then  $L$  is a residuated lattice.*

*Consider the L-VFSM,  $\mathcal{M} = (Q, \Sigma, \delta)$ , where  $Q = \{p_1, p_2, p_3, p_4\}$ ,  $\Sigma = \{\sigma\}$ , and  $\delta$  is represented by the following diagram:*



Let  $A = \{p_2, p_3, p_4\}$ . Then  $\mu_\sigma(A) = \{p_3, p_4\}$  and  $\mu_\sigma(\mu_\sigma(A)) = \{p_4\}$ . This shows that  $\mu_\sigma(\mu_\sigma(A)) \neq \mu_\sigma(A)$ . Also, if  $B = \{p_1, p_2, p_3\}$ , then  $\mu_s(B) = \{p_1, p_2\}$  and  $\mu_s(\mu_s(B)) = \{p_1\}$ . This shows that  $\mu_s(\mu_s(B)) \neq \mu_s(B)$ .

**Theorem 4.2.** Let  $\mathcal{M} = (Q, \Sigma, \delta)$  be an  $L$ -VFSM with  $L$  as complete residuated lattice without zero divisors. Then for all  $A \subseteq Q$ ,  $\mu_\sigma(\mu_\sigma(A)) = \mu_\sigma(A)$  and  $\mu_s(\mu_s(A)) = \mu_s(A)$ .

**Proof.** We prove only,  $\mu_\sigma(\mu_\sigma(A)) = \mu_\sigma(A)$ . Let  $q \in \mu_\sigma(A)$ . Then  $\sigma(A) \subseteq A$ . In order to show that  $q \in \mu_\sigma(\mu_\sigma(A))$ , it is enough to show that  $\sigma(q) \subseteq \mu_\sigma(A)$ . So let  $p \in \sigma(q)$ . Then  $\delta^*(p, x, q) > 0$  for some  $x \in \Sigma^*$ . Also, if  $r \in \sigma(p)$ , we have  $\delta^*(r, y, p) > 0$  for some  $y \in \Sigma^*$ . Since,  $L$  being without zero divisors,  $\delta^*(p, x, q) * \delta^*(r, y, p) > 0$ . Thus,  $\delta^*(r, yx, q) > 0$ , whereby  $r \in \sigma(q)$ . But as  $\sigma(q) \subseteq A$ ,  $r \in A$ . Thus  $\sigma(p) \subseteq A$ , implying that  $\sigma(q) \subseteq \mu_\sigma(A)$ . Also, as  $\mu_\sigma(\mu_\sigma(A)) \subseteq \mu_\sigma(A)$ , we have  $\mu_\sigma(\mu_\sigma(A)) = \mu_\sigma(A)$ .

**Definition 4.2.** [33] Let  $\mathcal{M} = (Q, \Sigma, \delta)$  be an  $L$ -VFSM. A subset  $R \subseteq Q$  is called

- (i) genetic if  $\sigma(R) \subseteq s(R)$ .
- (ii) genetically closed if  $\exists P \subseteq R$  such that  $\sigma(P) \subseteq s(P)$  and  $s(P) = R$ .

**Theorem 4.3.** Let  $\mathcal{M} = (Q, \Sigma, \delta)$  be an  $L$ -VFSM and  $R \subseteq Q$ . Then  $R$  is

- (i) genetic iff  $R \subseteq \mu_\sigma(s(R))$ .
- (ii) genetically closed iff  $\exists P \subseteq R$  such that  $P \subseteq \mu_\sigma(R)$  and  $s(P) = R$ .

**Proof.** (i) Let  $R$  be a genetic and  $r \in R$ . Then  $\sigma(r) \subseteq \sigma(R)$ . But as  $R$  is genetic, we have  $\sigma(R) \subseteq s(R) \Rightarrow \sigma(r) \subseteq s(R) \Rightarrow r \in \mu_\sigma(s(R))$  and hence  $R \subseteq \mu_\sigma(s(R))$ . Conversely, let  $R \subseteq \mu_\sigma(s(R))$  and  $q \in \sigma(R) \Rightarrow \exists p \in R$  such that  $q \in \sigma(p)$ . Since  $p \in R$ , we have  $p \in \mu_\sigma(s(R)) \Rightarrow \sigma(p) \subseteq s(R) \Rightarrow q \in s(R)$ . Thus  $\sigma(R) \subseteq s(R)$ . Hence,  $R$  is genetic.

(ii) Follows from (i) and the definition 4.2.



## 5. Fuzzy initial core and fuzzy final core operators

In this section we introduce two new fuzzy operators say fuzzy initial core and fuzzy final core. We show that these operators are the generalizations of source core and successor core operators respectively.

**Definition 5.1.** Let  $\mathcal{M} = (Q, \Sigma, \delta)$  be an L-VFSM and  $A \in \mathcal{F}_L(Q)$ . Then the fuzzy initial core and fuzzy final core of  $A$  are defined respectively as

$$\begin{aligned} Ico(A)(q) &= \bigwedge_{\substack{x \in \Sigma^* \\ p \in Q}} \{ \delta^*(p, x, q) \rightarrow A(p) \} \text{ and} \\ Fco(A)(q) &= \bigwedge_{\substack{x \in \Sigma^* \\ p \in Q}} \{ \delta^*(q, x, p) \rightarrow A(p) \} \end{aligned}$$

**Remark 5.1.** From the following theorem one can conclude that fuzzy initial core ( $Ico$ ) and fuzzy final core ( $Fco$ ) are the generalization of source core ( $\mu_\sigma$ ) and successor core ( $\mu_s$ ) respectively.

**Theorem 5.1.** Let  $\mathcal{M} = (Q, \Sigma, \delta)$  be an L-VFSM. Then for all  $q \in Q$ ,

- (i)  $Ico(1_A)(q) = 1$  iff  $q \in \mu_\sigma(A)$
- (ii)  $Fco(1_A)(q) = 1$  iff  $q \in \mu_s(A)$ .

**Proof.** (i) Let  $A \subseteq Q$  and  $q \in Q$ . Suppose  $Ico(1_A)(q) = 1$ . Then  $\bigwedge_{\substack{x \in \Sigma^* \\ p \in Q}} \{ \delta^*(p, x, q) \rightarrow 1_A(p) \} = 1 \Rightarrow$  for all  $p \in Q, x \in \Sigma^*, \delta^*(p, x, q) \rightarrow 1_A(p) = 1 \Rightarrow$  for all  $p \in Q, x \in \Sigma^*, \delta^*(p, x, q) \leq 1_A(p)$ . If,  $p \in \sigma(q)$ , we have  $\delta^*(p, x, q) > 0$ , for some  $x \in \Sigma^*$ . Therefore,  $1_A(p) > 0$  and thus  $p \in A$ . Hence,  $\sigma(q) \subseteq A \Rightarrow q \in \mu_\sigma(A)$ .

Conversely, suppose that  $q \in \mu_\sigma(A) \Rightarrow \sigma(q) \subseteq A$ . For, any  $p \in Q, x \in \Sigma^*$ , if  $\delta^*(p, x, q) = 0$  then  $\delta^*(p, x, q) \leq 1_A(p)$ , also if  $\delta^*(p, x, q) > 0$  then  $p \in \sigma(q) \subseteq A$ , there fore,  $1_A(p) = 1$  and thus  $\delta^*(p, x, q) \leq 1_A(p) \implies$  for all  $p \in Q, x \in \Sigma^*, \delta^*(p, x, q) \leq 1_A(p) \Rightarrow$  for all  $p \in Q, x \in \Sigma^*, \delta^*(p, x, q) \rightarrow 1_A(p) \} = 1 \Rightarrow \bigwedge_{\substack{x \in \Sigma^* \\ p \in Q}} \{ \delta^*(p, x, q) \rightarrow 1_A(p) \} = 1$  i.e.  $Ico(1_A)(q) = 1$ .

- (ii) On the same line we can prove.

We now depict few properties of these operators.

**Theorem 5.2.** Let  $\mathcal{M} = (Q, \Sigma, \delta)$  be an L-VFSM. Then for all  $A, B \in \mathcal{F}_L(Q)$ ,

i) if  $A \subseteq B$ , then  $Ico(A) \subseteq Ico(B)$  and  $Fco(A) \subseteq Fco(B)$

ii)  $Ico(A) \subseteq A$  and  $Fco(A) \subseteq A$

iii)  $Ico(A \cap B) = Ico(A) \cap Ico(B)$  and  $Fco(A \cap B) = Fco(A) \cap Fco(B)$

iv)  $Ico(Ico(A)) = Ico(A)$  and  $Fco(Fco(A)) = Fco(A)$ .

**Proof.** We only prove the properties of  $Ico$ , because the properties of  $Fco$  are similar.

(i) Let  $A \subseteq B$ .

Then  $A(p) \leq B(p), \forall p \in Q$ .

Now for any  $q \in Q$ , we have

$$Ico(A)(q) = \bigwedge_{\substack{x \in \Sigma^* \\ p \in Q}} \{\delta^*(p, x, q) \rightarrow A(p)\} \leq \bigwedge_{\substack{x \in \Sigma^* \\ p \in Q}} \{\delta^*(p, x, q) \rightarrow B(p)\} = Ico(B)(q).$$

Thus,  $Ico(A) \subseteq Ico(B)$ .

(ii) Let  $q \in Q$ .

$$\text{Then } Ico(A)(q) = \bigwedge_{\substack{x \in \Sigma^* \\ p \in Q}} \{\delta^*(p, x, q) \rightarrow A(p)\} \leq \delta^*(q, \varepsilon, q) \rightarrow A(q) = 1 \rightarrow A(q) = A(q).$$

Thus,  $Ico(A) \subseteq A$ .

(iii) Let  $q \in Q$ . Then

$$\begin{aligned} Ico(A \cap B)(q) &= \bigwedge_{\substack{x \in \Sigma^* \\ p \in Q}} \{\delta^*(p, x, q) \rightarrow (A \cap B)(p)\} \\ &= \bigwedge_{\substack{x \in \Sigma^* \\ p \in Q}} \{\delta^*(p, x, q) \rightarrow (A(p) \wedge B(p))\} \\ &= \bigwedge_{\substack{x \in \Sigma^* \\ p \in Q}} \{(\delta^*(p, x, q) \rightarrow A(p)) \wedge (\delta^*(p, x, q) \rightarrow B(p))\} \\ &= \bigwedge_{\substack{x \in \Sigma^* \\ p \in Q}} \{\delta^*(p, x, q) \rightarrow A(p)\} \wedge \bigwedge_{\substack{x \in \Sigma^* \\ p \in Q}} \{\delta^*(p, x, q) \rightarrow B(p)\} \\ &= Ico(A)(q) \wedge Ico(B)(q) \\ &= (Ico(A) \cap Ico(B))(q). \end{aligned}$$

Thus,  $Ico(A \cap B) = Ico(A) \cap Ico(B)$ .

(iv) Let  $q \in Q$ . Then

$$\begin{aligned}
 Ico(Ico(A))(q) &= \bigwedge_{\substack{x \in \Sigma^* \\ p \in Q}} \{\delta^*(p, x, q) \rightarrow Ico(A)(p)\} \\
 &= \bigwedge_{\substack{x \in \Sigma^* \\ p \in Q}} \{\delta^*(p, x, q) \rightarrow \bigwedge_{\substack{y \in \Sigma^* \\ r \in Q}} \{\delta^*(r, y, p) \rightarrow A(r)\}\} \\
 &= \bigwedge_{\substack{x \in \Sigma^* y \in \Sigma^* \\ p \in Q r \in Q}} \{\delta^*(p, x, q) \rightarrow (\delta^*(r, y, p) \rightarrow A(r))\} \\
 &= \bigwedge_{\substack{x \in \Sigma^* y \in \Sigma^* \\ p \in Q r \in Q}} \{(\delta^*(p, x, q) * \delta^*(r, y, p)) \rightarrow A(r)\} \\
 &= \bigwedge_{\substack{x, y \in \Sigma^* \\ r \in Q}} \{\bigvee_{p \in Q} (\delta^*(p, x, q) * \delta^*(r, y, p)) \rightarrow A(r)\} \\
 &= \bigwedge_{\substack{x, y \in \Sigma^* \\ r \in Q}} \{\delta^*(r, yx, q) \rightarrow A(r)\} = Ico(A)(q).
 \end{aligned}$$

Thus,  $Ico(Ico(A)) = Ico(A)$ .

**Theorem 5.3.** Let  $\mathcal{M} = (Q, \Sigma, \delta)$  be an L-VFSM. Then for any  $A \in \mathcal{F}_L(Q)$ ,

(i)  $so(A) = A$  iff  $Ico(A) = A$ ,

(ii)  $su(A) = A$  iff  $Fco(A) = A$ .

**Proof.** (i) Let  $so(A) = A$ . For any  $q \in Q$ ,  $so(A)(q) = A(q) \Rightarrow$  for any  $q \in Q$ ,  $\bigvee_{\substack{x \in \Sigma^* \\ p \in Q}} \{\delta^*(q, x, p) * A(p)\} = A(q) \Rightarrow$  for any  $p, q \in Q, x \in \Sigma^*$ ,  $\delta^*(q, x, p) * A(p) \leq A(q) \Rightarrow$  for any  $p, q \in Q, x \in \Sigma^*$ ,  $A(p) \leq \delta^*(q, x, p) \rightarrow A(q) \Rightarrow$  for any  $p \in Q$ ,  $A(p) \leq \bigwedge_{\substack{x \in \Sigma^* \\ p \in Q}} \{\delta^*(q, x, p) \rightarrow A(q)\} \Rightarrow$  for any  $p \in Q$ ,  $A(p) \leq Ico(A)(p)$  whence,  $A \leq Ico(A)$ . But, by Theorem 5.2  $Ico(A) \leq A$  and thus  $Ico(A) = A$ . Similarly, we can prove converse.

(ii) On the same line we can prove the result.

## 6. Conclusion

In this paper we have introduced successor core operator in an analogue to the source core operator. Two more new fuzzy operators namely fuzzy initial core and fuzzy final core are introduced and their properties are studied.

## Conflict of Interests

The authors declare that there is no conflict of interests.

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